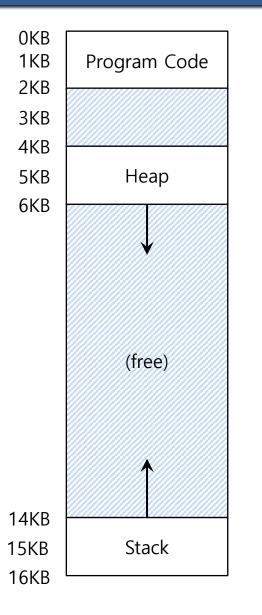
16. Segmentation

Operating System: Three Easy Pieces

Inefficiency of the Base and Bound Approach



- **Big chunk of "free"** space
- "free" space **takes up** physical memory.
- Hard to run when an address space does not fit into physical memory

- Segment is just a contiguous portion of the address space of a particular length.
 - Logically-different segment: code, stack, heap
- **•** Each segment can be **placed** in **different part of physical memory**.
 - Base and bounds exist per each segment.

Placing Segment In Physical Memory

0KB **Operating System** 16KB (not in use) Segment Base Stack Code (not in use) Heap 32KB Code Stack Heap 48KB (not in use) MMU 64KB **Physical Memory**

Size

2K

2K

2K

32K

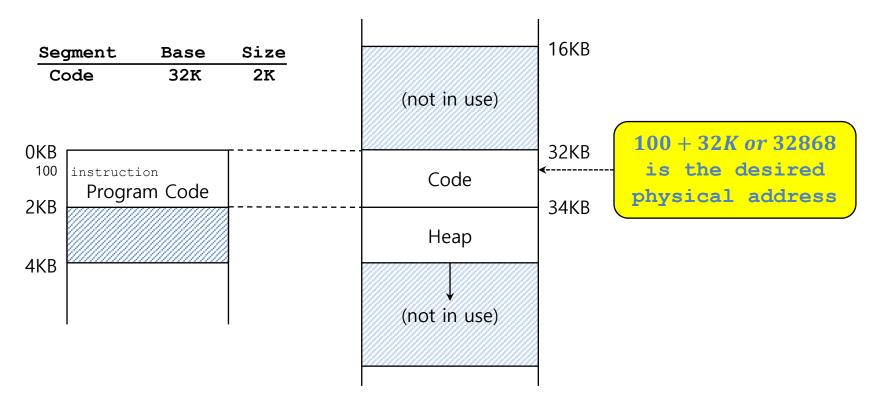
34K

28K

Address Translation on Segmentation

physical address = offset + base

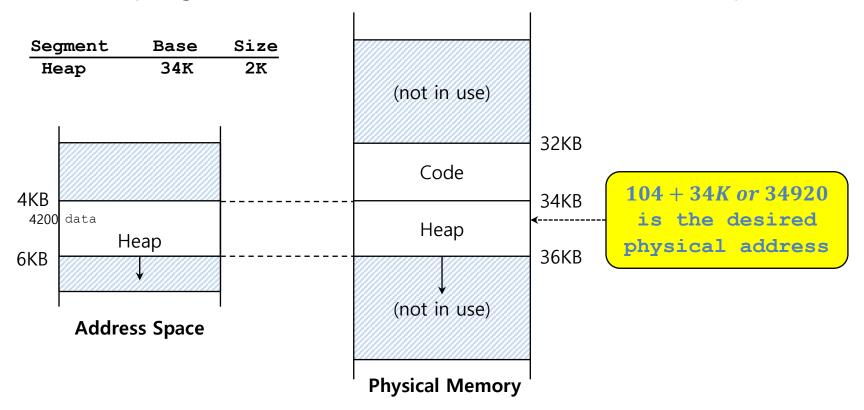
- **The** offset of virtual address 100 is 100.
 - The code segment **starts at virtual address 0** in address space.



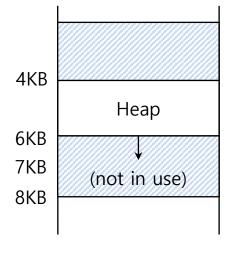
Virtual address + base is not the correct physical address!

The offset of virtual address 4200 is 104.

• The heap segment starts at virtual address 4096 in address space.



- If an illegal address such as 7KB which is beyond the end of heap is referenced, the OS occurs segmentation fault.
 - The hardware detects that address is **out of bounds**.

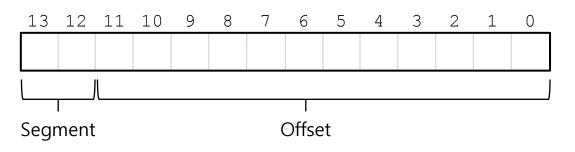


Address Space

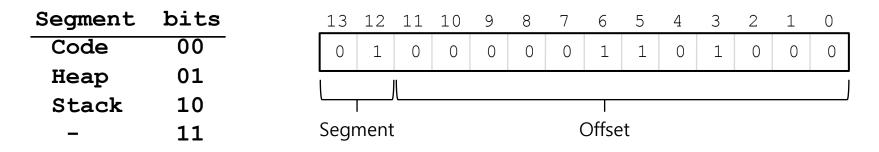
Referring to Segment

Explicit approach

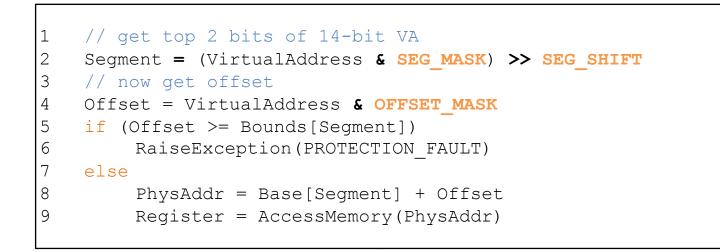
 Chop up the address space into segments based on the top few bits of virtual address.



Example: virtual address 4200 (01000001101000)

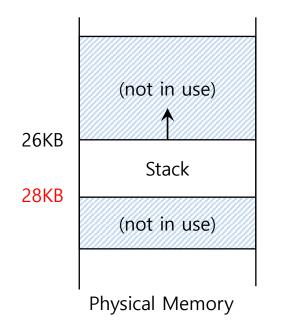


Referring to Segment(Cont.)



- SEG_MASK = 0x3000(110000000000)
- SEG SHIFT = 12
- OFFSET MASK = 0xFFF (001111111111)

- **D** Stack grows **backward**.
- **Extra hardware support** is need.
 - The hardware checks which way the segment grows.
 - 1: positive direction, 0: negative direction



Segment Register(with Negative-Growth Support)

Segment	Base	Size	Grows Positive?
Code	32K	2K	1
Heap	34K	2K	1
Stack	28K	2K	0

Support for Sharing

- **D** Segment can be **shared between address** space.
 - Code sharing is still in use in systems today.
 - by extra hardware support.
- **D** Extra hardware support is need for form of **Protection bits**.
 - A few more bits per segment to indicate permissions of read, write and execute.

Segment	Base	Size	Grows Positive?	Protection
Code	32K	2K	1	Read-Execute
Heap	34K	2K	1	Read-Write
Stack	28K	2K	0	Read-Write

Segment Register Values(with Protection)

Fine-Grained and Coarse-Grained

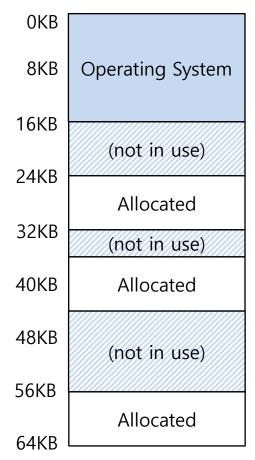
- **Coarse-Grained** means segmentation in a small number.
 - e.g., code, heap, stack.
- Fine-Grained segmentation allows more flexibility for address space in some early system.
 - To support many segments, Hardware support with a segment table (stored in memory) is required.
 - Exploit temporal locality (mem. usage)

OS support: Fragmentation

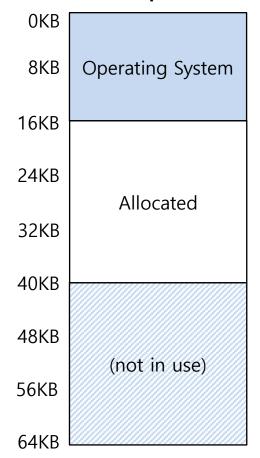
- External Fragmentation: little holes of free space in physical memory that make difficulty to allocate new segments.
 - There is **24KB free**, but **not in one contiguous** segment.
 - The OS cannot satisfy the **20KB request**.
- **Compaction**: **rearranging** the exiting segments in physical memory.
 - Compaction is **costly**.
 - **Stop** running process.
 - Copy data to somewhere.
 - Change segment register value.
- 1000 ways to solve it
 - None of them are the "best"
- Added to creating, terminating, and context switches

Memory Compaction





Compacted



This lecture slide set is used in AOS course at University of Cantabria. Was initially developed for Operating System course in Computer Science Dept. at Hanyang University. This lecture slide set is for OSTEP book written by Remzi and Andrea Arpaci-Dusseau (at University of Wisconsin)